FPGA Implementation of CORDIC Multiplier based Redundant Floating-point Butterfly Architecture for Signal Processing Applications

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Abstract: The main objective of this paper is to design a area efficient and high speed floating-point (FP) butterfly architecture for a signal processing applications. The butterfly architectures used in Fast Fourier Transform (FFT) algorithms. Most of the butterfly architectures uses fixed point arithmetic than the FP arithmetic because of lesser speed but main advantage is the high dynamic range. The limitation of FP arithmetic is overcome by using a FDPA unit but it consume more area. This draw back overcome by using a Coordinate Rotation Digital Computer (CORDIC) algorithm. In this paper we propose a CORDIC Multiplier based redundant floating point butterfly architecture for a signal processing applications. The proposed work synthesis is carried out by using XILINX ISE synthesis tool on the XILINX 14.7 platform and simulated using Model Sim simulator. The performance of this design is evaluated on XILINX Spartan3 family device.

Keywords: CORDIC algorithm, FDPA unit, Floating-point, Fixed –point

I. INTRODUCTION

FFT algorithms use butterfly architectures to compute the Discrete Fourier Transform (DFT) of input signal. The butterfly circuit consists of several multipliers and an adder over complex number [1] is shown in Fig.1. Most of the butterfly architectures have been using fixed-point arithmetic, until recently that butterfly architectures based on FP operations grown [1]. The main advantage of FP arithmetic is the high dynamic range, but at the expense of higher delays [1].
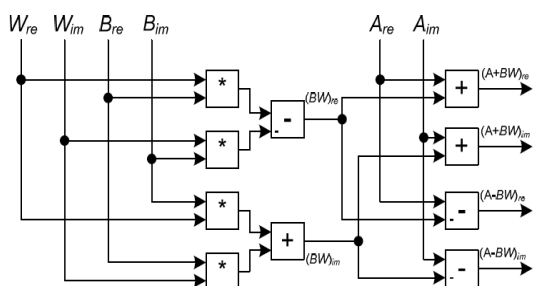

Fig.1. DIT Butterfly architecture expanded with complex numbers

As Discussed, The main limitation of the FP operation is their high delay in comparison with the fixed − point counter parts, is as a way to reducing delay of the FP arithmetic is to merge the several operations in a single floating − point unit. Here are the few techniques for merging the operations. They are Fused-Multiply-Add, Dot-Product and Fused Dot-Product-Add.
The time and area will be saved for butterfly architecture [2][5] by using the FP Fused-Multiply-Add technique, but the butterfly architecture cannot be implemented directly. In order to circumvent this problem Dot-Product unit is

required. The Dot-Product unit is an extension of Fused-Multiply technic.
The Dot-Product unit calculates $AB + CD$ (or) $AB − CD$. The Dot-Product unit eliminates more intermediate Normalization and Rounding operations because of it combines more FP operations hence saving more area and time[3] than Fused-Multiply-Add.
Fused-Dot-Product-Add(FDPA) unit is the extension to the Dot-Product unit. By this unit combining more FP operations than the Dot-Product unit. FDPA unit calculates $AB \pm CD \pm E$ over FP operands [1]. The butterfly unit implemented using the FDPA unit is as follows. By using FDPA units achieves higher speed but area consumption is more. In order to achieve higher speed with lesser area CORDIC algorithm has been introduced in FDPA unit [1].
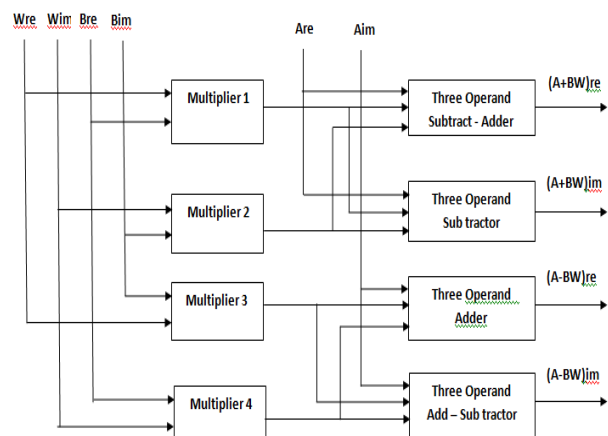

Fig.2 Butterfly architecture with FP Fused-Dot-Product-Add

In this paper CORDIC Multilier based butterfly architecture is proposed which has higher speed and lesser area. The paper structured as follows Section II explains the Proposed Butterfly Architecture, Section III Explains Experimental Results of Proposed Butterfly Architecture, Section IV Conclusion.

II. PROPOSED BUTTERFLY ARCHITECTURE

The proposed CORDIC Multiplier based butterfly architecture is shown in below Fig.3. Proposed butterfly architecture consists of CORDIC Multiplier, Three Operand Subtract-Adder, Three Operand Subtractor, Three Operand Adder and Three Operand Add-Subtractor.
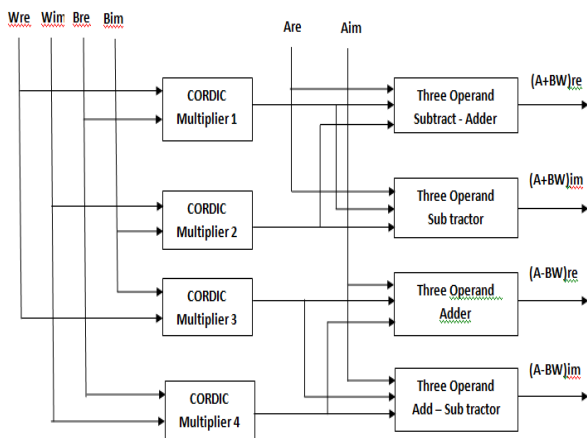

Fig.3. Proposed Butterfly Architecture

A. CORDIC Multiplier

The CORDIC algorithm uses a one-bit-at-a-time approach to make computations to an arbitrary precision. In the process, relatively small lookup tables are used for constants necessary for the algorithm. Typically these tables require only one to two entries per bit of precision. CORDIC algorithms also use only right shifts and additions, minimizing the computation time [10].

A CORDIC algorithm for Multiplication may be derived by using a series representation for a as follows:

$$c = a * b \qquad (1)$$

$$= b * \sum_{i=1}^{B} x_i * 2^{-i} \qquad (2)$$

$$= \sum_{i=1}^{B} b * x_i * 2^{-i} \qquad (3)$$

$$= \sum_{i=1}^{B} x_i * (b * 2^{-i}) \qquad (4)$$

From this it is seen that c is composed of shifted versions of b. The unknown coefficients, x_i, may be found by driving a to zero one bit at a time. If the ith bit of a is non-zero, b_i is right shifted by i bits and added to the current value of c. The ith bit is then removed from a by subtracting 2^{-i} from a. If a is negative, the ith bit in the twos complement format would be removed by adding 2^{-i}. In

either case, when a has been driven to zero all bits have been examined and c contains the signed product of a and b correct to B bits.

This algorithm is similar to the standard shift and add multiplication algorithm except for two important features. First, arithmetic right shifts are used instead of left shifts, allowing signed numbers to be used. Secondly, computing the product to B bits with the CORDIC algorithm is equivalent to rounding the result of the standard algorithm to the most significant B bits. The final algorithm is as follows:

```
Multiply (a. b) {
    for ( j = 1 ; j =< R ; j + +) {
        if( a > 0)
            a = a – 2^−j
        c = c + b ∗ 2^−j
            else
            a = a + 2^−j
        c = c + b ∗ 2^−j
        }
    return(c)
    }
```

By using CORDIC Multiplication algorithm perform floating-point Multiplication on both signed and unsigned numbers.
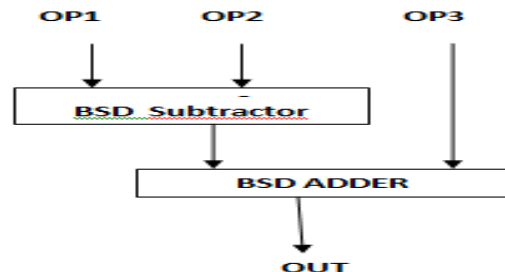
B. Three Operand Subtract- Adder

Fig.4. Block Diagram of Three Operand Subtract- Adder

Three Operand Subtract-Adder consists of Binary Signed Digit (BSD) subtractor and BSD adder. These are cascaded to form Three Operand Subtract-Adder is shown in Fig.4. It has three inputs, two are from multiplier 1 and 2 and the real part of A (Are) will be given as third input as shown in the Fig.3.
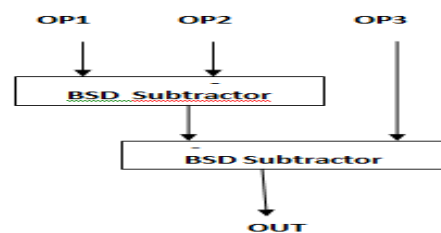
C. Three Operand Subtractor

Fig.5 Block Diagram of Three Operand Subtractor

Three Operand Subtractor consists of two BSD subtractors. These are cascaded to form Three Operand Subtractor is shown in Fig.5. It has three inputs, two are from multiplier 1 and 2 and the real part of A (Are) will be given as third input as shown in the Fig.3.
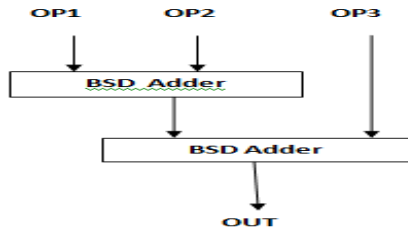
D. Three Operand Adder

Fig.6. Block Diagram of Three Operand Adder.

Three Operand Adder consists of two BSD Adders . These are cascaded to form Three Operand Adder is shown in Fig.6. It has three inputs, two are from multiplier 3 and 4 and the imaginary part of A (Aim) will be given as third input as shown in the Fig.3.

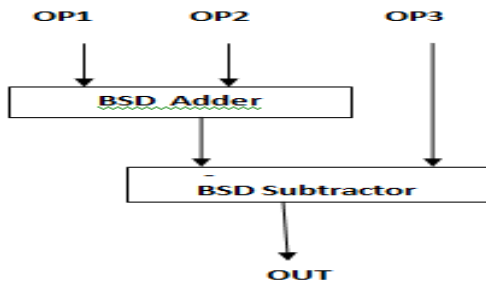E. Three Operand Add-Subtractor

Fig.7. Block Diagram of Three Operand Add-Subtractor.

Three Operand Subtract-Adder consists of BSD Adder and BSD Subtractor. These are cascaded to form Three Operand Add-Subtractor is shown in Fig.7. It has three inputs, two are from multiplier 3 and 4 and the Imaginary part of A (Aim) will be given as third input as shown in the Fig.3.

The proposed CORDIC Multiplier based butterfly architecture is designed in Xilinx ISE and there RTL schematic diagram is shown in Fig 8 and Fig 9.

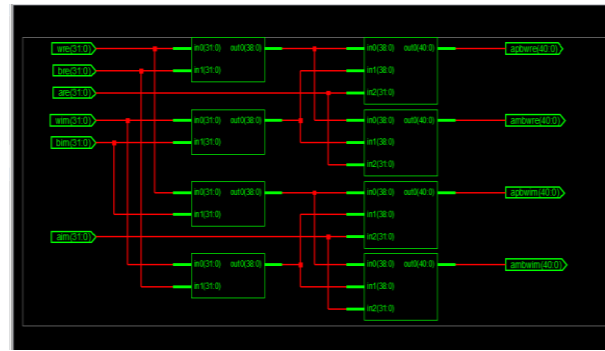Fig.8. RTL Schematic Diagram of Proposed Butterfly Architecture

Fig.9. Internal RTL Schematic Diagram of Proposed Butterfly Architecture.

III.EXPERIMENTAL RESULTS

The proposed butterfly architecture is simulated using ModelSim and the results are verified by using MATLAB. The CORDIC Multiplier simulation results are obtained by giving the following data as input In0 = 00000001010000000000000000000000 (1.25) and In2 = 00000001100000000000000000000000 (1.5) and obtain the output as out= 00000001100000000000000000000000 (1.5) is shown in Fig.10 and Fig.11.
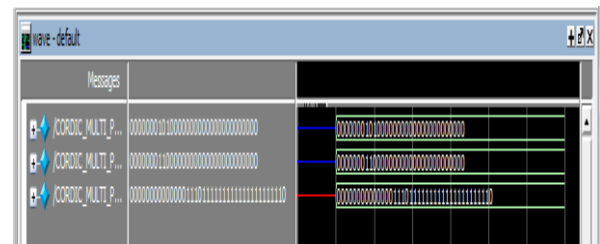

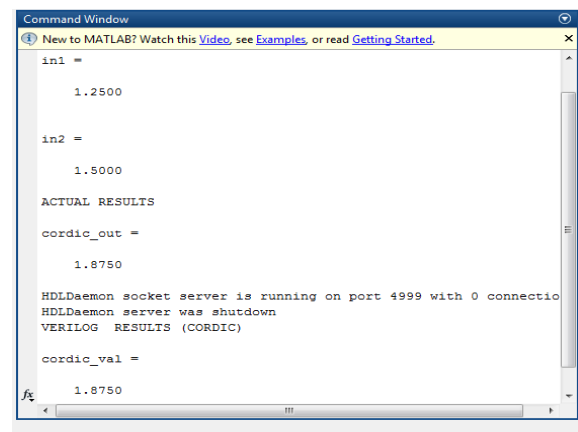Fig.10. Simulation Results of CORDIC Multiplier

Fig.11. Simulation results of CORDIC Multiplier verification using MATLAB.

The simulation results of Three Operand subtract-adder obtained by giving input In0= 00000000000001100001000111000100 0011001 (3.0347), In1 = 00000000000000110110010010101000 1100001 (1.6966), In2 = 00000000001101100111001110000010 (0.2127) and obtain output out = 00000000000000011000110100000001 00111010 (1.5508) is shown in Fig 12 and Fig 13.
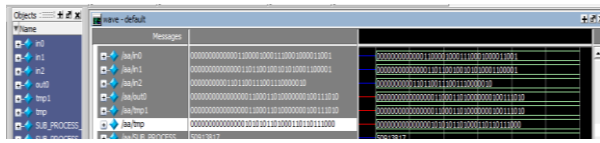
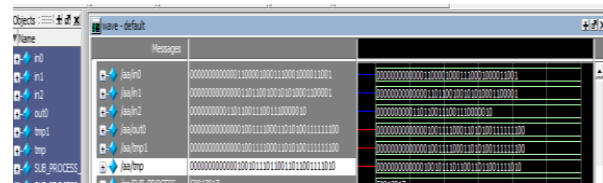Fig.12. Simulation Results of Three Operand Subtract – Adder

Fig.13. Simulation results of Three Operand Subtract-Adder verification using MATLAB.
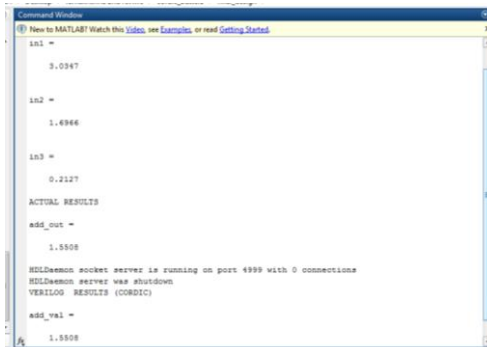
The simulation results of Three Operand subtractor obtained by giving input In0= 000000000000001100001000111000100001001 (3.0347), In1 = 0000000000000001101100100101010001100001 (1.6966), In2 = 00000000001101100111001110000010 (0.2127) and obtain output out = 1111111111111111011011111111100010111001010 (-1.1254) is shown in Fig. 14 and Fig. 15.

Fig.14. Simulation Results of Three Operand Subtractor

Fig.15. Simulation results of Three Operand Subtractor verification using MATLAB.

The simulation results of Three Operand adder obtained by giving input In0= 000000000000001100001000111000100001001 (3.0347), In1 = 0000000000000001101100100101010001100001 (1.6966), In2 = 00000000001101100111001110000010 (0.2127) and obtain output out = 0000000000000001001111000110101001111111100 (4.9440) is shown in Fig 16 and Fig 17.
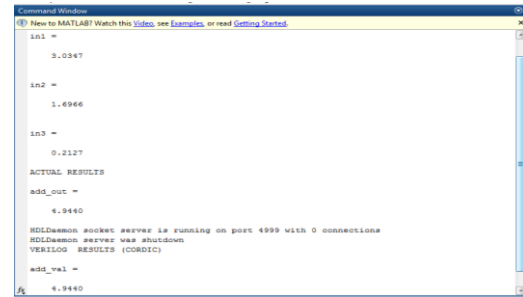

Fig.16. Simulation Results of Three Operand Adder

Fig.17. Simulation results of Three Operand Adder verification using MATLAB.

The simulation results of Three Operand Add-Subtractor obtained by giving input In0= 000000000000001100001000111000100001001 (3.0347), In1 = 0000000000000001101100100101010001100001 (1.6966), In2 = 00000000001101100111001110000010 (0.2127) and obtain output out = 1111111111111110110111101100111101000011000 (-4.5186) is shown in Fig 18 and Fig 19.
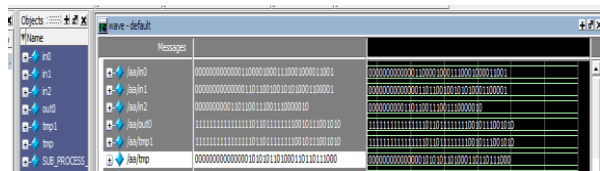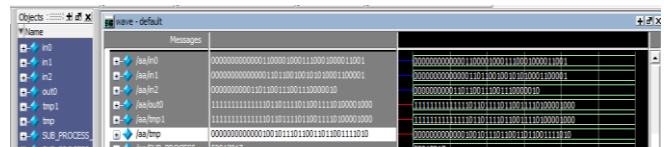

Fig.18. Simulation Results of Three Operand Add-Subtractor

Fig.19. Simulation results of Three Operand Add-Subtractor verification using MATLAB.

The simulation results of proposed butterfly architecture obtained by giving inputs are

are= 0000000110011001110110110010011 (1.6010),
aim = 00000001010000000000000000000000 (1.2500),
bre = 00000001100000000000000000000000 (1.5000),
bim = 00000010100000000000000000000000 (2.5000),
wre = 000000001011010011111110111110100 (0.7070),
wim = 000000001011010011111110111110100 (0.7070).
and obtain outputs are
apbwre=000000000000000000011100100110111010010111 1 (0.8940),

apbwim = 0000000000000010000010011111011111001100 (4.0780),
ambwre = 0000000000000001001001110110110010001011 (2.3080)
ambwim = 1111111111111111001101100000010000110100 (-1.5780) is shown in Fig 20 and Fig 21 and Fig 22.
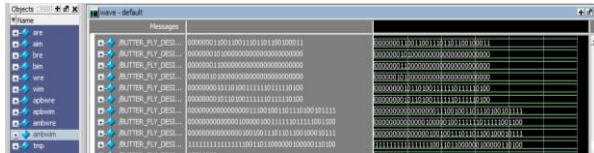

Fig.20. Simulation Results of Proposed Butterfly Architecture

Fig.21. Simulation results input values of Proposed Butterfly Architecture verification using MATLAB

Fig.22. Simulation result output values of Proposed Butterfly Architecture verification using MATLAB

Table I. Comparision of FP Butterfly Architectures

Parameters	Existing work	Proposed work
Number of Slice LUTs	2168	1044
Number of 4 input LUTs	2576	2024
Number of Bonded IOBs	944	310
Time Delay	60.4ns	53.296ns

IV CONCLUSION

We designed a high speed and area efficient Floating-Point butterfly architecture, which is faster than existing works. The reason for getting lesser area is Multiplier designed by using CORDIC algorithm which reduces the complexity, and reason for speed improvement is FDPA unit.

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